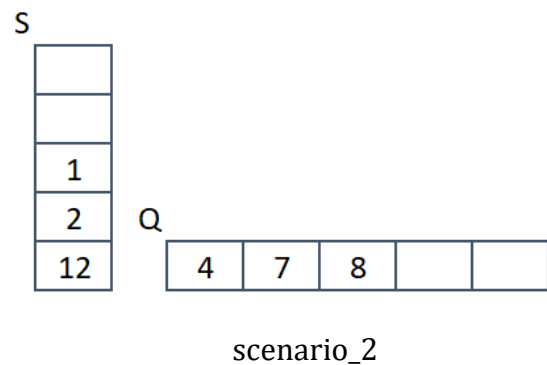
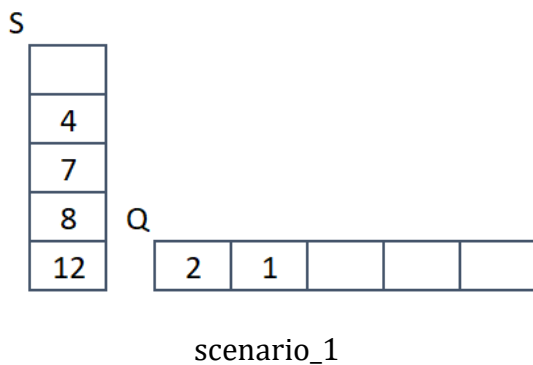


Solve all Questions

Q1 [2 points] Give the Big-Oh notation for the following functions

$f(N) = N^2 + \log N^2 + 2N \log N$	$O(N^2)$
$f(N) = (N \cdot (100N + 5000000))^2$	$O(N^4)$
$f(N) = N^{1/2} + \log(\log N)$	$O(N)$
$f(N) = 1000^{100} + \log N$	$O(\log N)$

Q2. [2 points] Assume you have a non-empty Stack S and a Queue Q of type integers. Provide a sequence of calls to modify the contents of S and Q from scenario\_1 to scenario\_2.



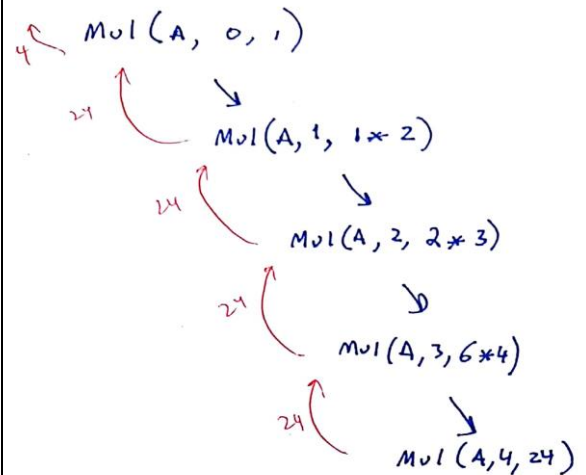
- Q.enqueue(S.pop())**
- Q.enqueue(S.pop())**
- Q.enqueue(S.pop())**
- S.push(Q.dequeue())**
- S.push(Q.dequeue())**

Q3 [4]. Given the following code, estimate the number of operations and describe the worst case running time in Big-Oh notation in terms of the variable n. Show trace where appropriate.

```
public int Mul(int [] A, int i, int t)
{
    if (i >= A.length-1)
        return t;
    else
        return Mul(A, ++i, t * A[i]);
}
```

**Mul() is recursively called i times. I would run up to n, which is the length of this array. The run time will be O(n)**

```
Mul(new int[] {1,2,3,4}, 0, 1)
```



```
public void diag(int n) {
    int count = 0;
    for (int i = 0; i < n; i++)
        for (int j = 0; j < i; j++)
            if(i==j)
                count++;
}
```

**The if statement will execute  $n/2 * (n+1)$  times**

$$1+2+3+4+5+\dots = \frac{n(n+1)}{2}$$

**which is  $O(n^2)$**

```
diag(5)
```

```
i:0 (1 times)
j=0
```

```
i:1 (2 times)
j=0
j=1
```

```
i:2 (3 times)
j=0
j=1
j=2
```

```
i:3 (4 times)
j=0
j=1
j=2
j=3
```

```
i:4 (5 times)
j=0
j=1
j=2
j=3
j=4
```

Q4. [3 points] Write an iterative method `public static double addDiagonal(double [][] D)` that adds and returns all the left-diagonal values of a  $n \times n$  2-dimensional array.

Show a trace to determine the runtime function  $T(n)$  estimating the number of operations.


```
public static double addDiagonal(double [][] D)
{
    double total = 0;
    for(int i=0;i<D.length;i++)
        total+=A[i][i];
    return total;
}
```

We must note that the size of the matrix is  $n \times n$ , which means it is a square matrix. We estimate the number of operation to be:

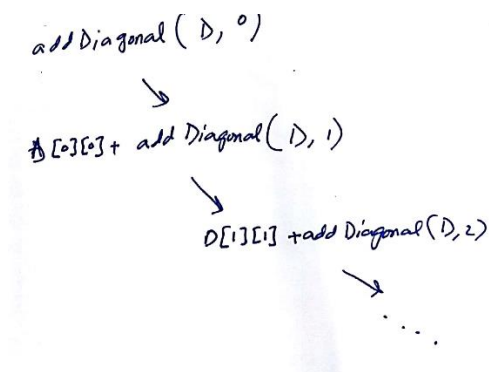
$$T(n) = 5n + 3$$

which is  $O(n)$

[+1 Bonus] Write the above method as a recursive method. Estimate the runtime of this method.

```
public static double addDiagonal(double [][] D, int i)
{
    if(i>=D.length)
        return 0;
    else
        return D[i][i] + addDiagonal(D, ++i);
}
```

We can estimate that the run time  $T(n) = 5c n$  which is  $O(n)$



Q5. [4 points] Modify the method insert that you defined in the double linked list class DList. The modified method has the following definition.

**public static void insert2(Node N, int pos)**

This method inserts a node *N* at position *pos* in the DList. You may start counting at the head. Make sure your method checks if value of *pos* is valid (i.e. *pos* must be a non-negative integer with value less than size).

```
public static void insert2(Node N, int pos){
    //we assume the first position is zero
    if(pos<0 || pos > size+1)
        return;
    if(pos==0)//add to the beginning
    {
        N.next = Head;
        N.prev = null;
        Head.prev = N;
        Head = N;
    }
    else if(pos == size) // we add at the end
    {
        N.next = null;
        N.prev = Tail;
        Tail.next =N;
        Tail = N;
    }
    else
    {
        //we add in the middle
        Node temp = Head;
        for(int i=1;i<pos;i++)
            temp = temp.next;

        N.next = temp.next;
        N.prev = temp.prev;
        temp.next = N;
        N.next.prev = N;
    }
    size++;
    return;
}
```