

## CS210 Data Structures (221) Final Exam

Name: \_\_\_\_\_ ID \_\_\_\_\_

Check your section:

<input type="checkbox"/> Dr. Basit Qureshi <input type="checkbox"/> Dr. Syed Umar Amin <input type="checkbox"/> Dr. Adel Ammar <input type="checkbox"/> Dr. Abdullah Alrajeh <input type="checkbox"/> Dr. Sajid Shah	<input type="checkbox"/> Dr. Najla Thuniyan <input type="checkbox"/> Dr. Sawsan Halawani <input type="checkbox"/> Dr. Alaa Shamasneh
--	--

Instructions:

- This exam contains four questions with multiple parts, on 6 sheets of papers (2 sided).
- Time allowed: 180 minutes
- Closed Book, Closed Notes.
- Use of Calculators is ALLOWED. Use of other computing devices / smartphones etc is strictly prohibited.
- Answer the problems on the exam sheets only. No additional attachments would be accepted.
- When the “time is over” is called, it is the students’ responsibility to submit his exam to the invigilator. Submitting completed exam 3 minutes after the “time is over” will incur a penalty of **5 points**.

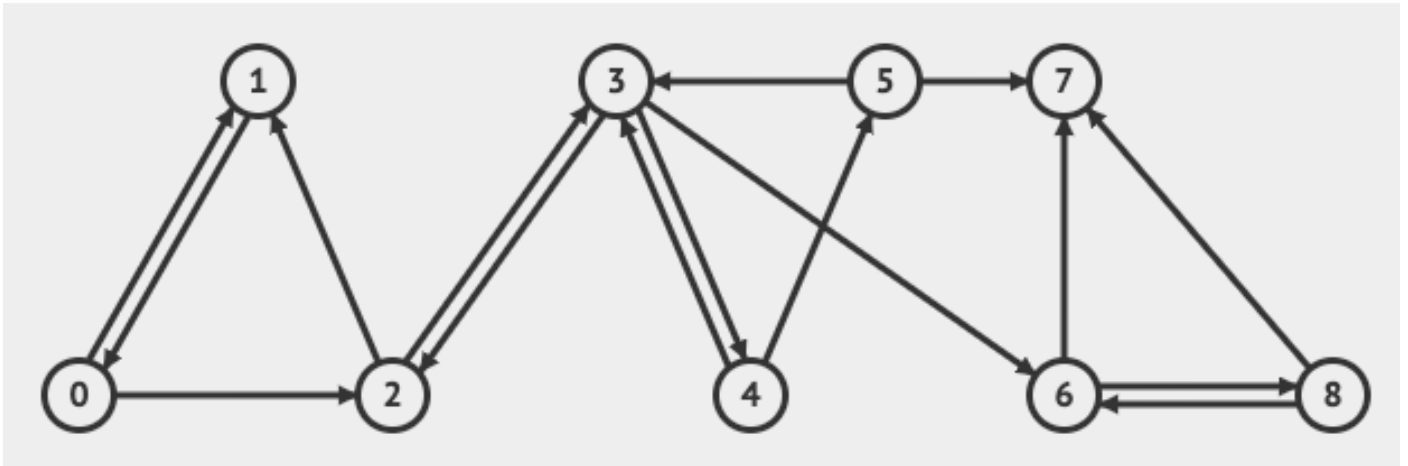
Few gentle reminders:

- If you get stuck on some problem for a long time, move on to the next one.
- You should be better off by first reading all questions and answering them in the order of what you think is the easiest to the hardest problem.
- Keep the points distribution in mind when deciding how much time to spend on each problem.

Question No.	Part a	Part b	Part c	Part d	Part e	Student’s Score
Question 1 (CLO 1)	/4	/3	/3	/3	/3	/16
Question 2 (CLO 2)	/2	/2	/4			/8
Question 3 (CLO 3)	/4	/4				/8
Question 4 (CLO 4)	/4	/4				/8
Total						/40

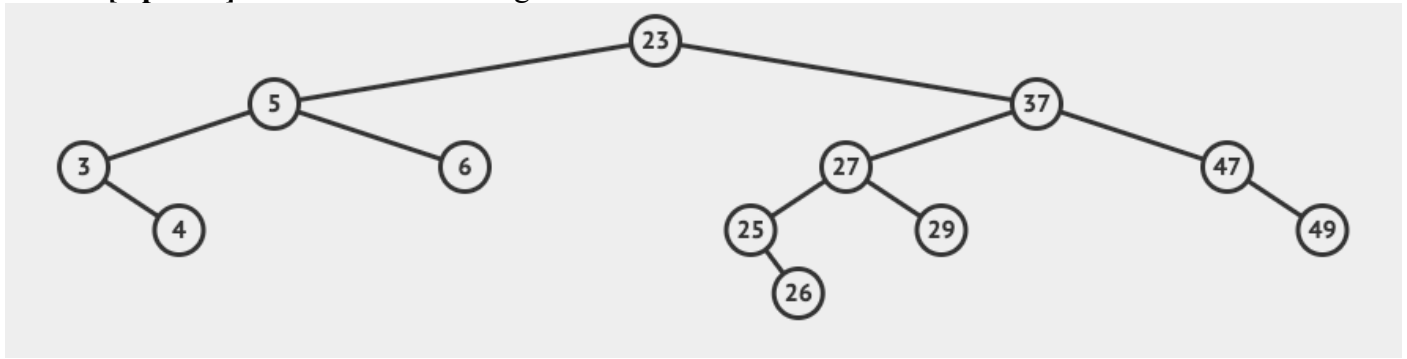


Question 1. [4 + 3 + 3 + 3 + 3 = 16 points] [CLO 1]



Part a. [2 + 1 + 1 points] Show the adjacency list for the directed graph illustrated above. Do a Depth First Search (DFS) and Breadth First Search (BFS) starting at 0. Choose the smaller value on vertices to select the next destination/hop.

**Part b. [3 points]** Observe the following AVL Tree.



**Remove 47** from the above tree. Re-Draw the tree and show appropriate rotations.

From the resulting tree you drew, **remove 6**. Re-Draw the tree and show appropriate rotations

From the resulting tree you drew above, **insert 28**. Re-Draw the tree and show appropriate rotations

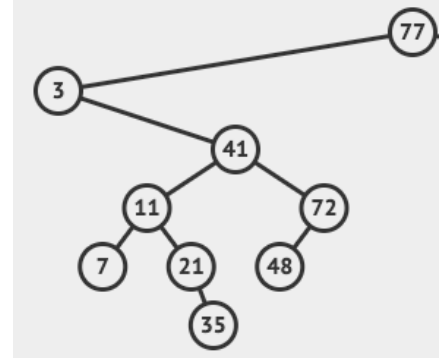
**Part c. [3 points]** Draw the max-heap given in the following array Show the heap operations for removing two values from this max-heap. Re-draw the heap after each removal.

0	1	2	3	4	5	6	7	8	9	10
<b>28</b>	<b>21</b>	<b>23</b>	<b>19</b>	<b>18</b>	<b>16</b>	<b>12</b>	<b>8</b>	<b>10</b>	<b>6</b>	<b>1</b>

**Part d. [3 points]** An AVL Tree is essentially a balanced BST. Write a method `Node Ancestors(int x)` that takes a parameter `int x`; this method finds the node N with value x. It returns a String containing all the ancestors of node N separated by spaces. Assume the tree stores integer *Key* in each node.

```
String Ancestors(int x) {
```

Example:



A call `Ancestors(21)` for this BST would return “11, 41, 3, 77”

**Part e. [3 points]** Write a method called **OddsToStack** that takes as parameter, a singly-linked list L, which contains an integer value in each node. It returns a stack that contains all the elements of the list except even numbers. Note the order of items is reversed!



```
public Stack OddsToStack(SinglyLinkedList L){
```

**Question 2. [2 + 2 + 4 = 8 points] [CLO 2]**

**Part a. [2 points]** The following methods search for an integer *key* in a 3D Array and returns true if found. Give their run-time.

<pre>public int search1 (int [][][] a, int key){     int p = a.length();     int r = a[0].length();     int c = a[0][0].length();     boolean FLAG = false;     for(int i=0;i&lt;p;i++){         for(int j=0;j&lt;r;j++){             for(int k=0;k&lt;c;k++){                 if(a[i][j][k]==key){                     FLAG = true;                 }             }         }     }     return FLAG; }</pre>	<pre>public int search2 (int [][][] a, int key){     int p = a.length();     int r = a[0].length();     int c = a[0][0].length();      for(int i=0;i&lt;p;i++){         for(int j=0;j&lt;r;j++){             for(int k=0;k&lt;c;k++){                 if(a[i][j][k]==key){                     return true;                 }             }         }     }     return false; }</pre>

Both programs are identical, which is better and why?

**Part b. [2 points]** Consider the following recursive method **mul** that takes two parameters int start and int end. The method multiplies all integer values between start and end returning the result as an integer value. Describe the worst-case running time of the function **mul**. Show/draw the recursion trace as necessary for this call.

**CALL: mul (2, 6)**

```
public static int mul(int start, int end) {
    if (start > end)
        return 1;
    else
        return end * mul(start, end -1)
}
```

**Part c. [4 points]** Identify the correct choice from the following MCQs:

	Choice?
<p>The runtime cost of a double left/right rotation in an AVL tree is</p> <p>a. <math>O(1)</math>  b. <math>O(\log n)</math>  c. <math>O(n)</math>  d. <math>O(n \log n)</math></p>	a
<p>The runtime cost of deleting a key from the AVL tree is</p> <p>a. <math>O(1)</math>  b. <math>O(\log n)</math>  c. <math>O(n)</math>  d. <math>O(n \log n)</math></p>	b
<p>Space required for storing an Un-directed Graph in an Adjacency Matrix is</p> <p>a. <math>O(1)</math>  b. <math>O(N)</math>  c. <math>O(V * E)</math>  d. <math>O(V * V)</math></p>	d
<p>Removing an edge from a given edge-list for a directed graph take this runtime:</p> <p>a. <math>O(E)</math>  b. <math>O(E * E)</math>  c. <math>O(V * E)</math>  d. <math>O(V * V)</math></p>	a
<p>The runtime cost of traversing and printing all node values in an AVL tree is</p> <p>a. <math>O(1)</math>  b. <math>O(\log n)</math>  c. <math>O(n)</math>  d. <math>O(n \log n)</math></p>	c
<p>Worst case scenario for searching the smallest value in a AVL Tree</p> <p>a. <math>O(n^2)</math>  b. <math>O(\log n)</math>  c. <math>O(n)</math>  d. <math>O(n \log n)</math></p>	b
<p>Searching a value in a min-heap</p> <p>a. <math>O(1)</math>  b. <math>O(\log n)</math>  c. <math>O(n)</math>  d. <math>O(n \log n)</math></p>	c
<p>Removing a node at the Tail of a doubly linked list</p> <p>a. <math>O(1)</math>  b. <math>O(\log n)</math>  c. <math>O(n)</math>  d. <math>O(n \log n)</math></p>	a



**Question 3. [6 + 2 = 6 points] [CLO 3]**

**Part a. [6 points]** Suppose you have to store the following values in a hash table, implemented using linear probing. The hash function used was the identity function,  $h(x) = x \bmod 13$ . Assume that the hash table has a size is of 10; insert the following in this hash-table.

**9, 27, 39, 16, 22, 35, 18, 8, 20, 28.**

0	1	2	3	4	5	6	7	8	9

Identify, how many collisions \_\_\_\_\_ and displacements \_\_\_\_\_ occurred.

The table size is 10. If we re-size it to 13, would it make any difference?

0	1	2	3	4	5	6	7	8	9	10	11	12

Identify, how many collisions \_\_\_\_\_ and displacements \_\_\_\_\_ occurred.

Assuming the same data is inserted in a Hash-table size 13, using separate chaining as a collision resolution method, Draw the Hash table.

0
1
2
3
4
5
6
7
8
9
10
11
12

Identify, how many collisions \_\_\_\_\_ and displacements \_\_\_\_\_ occurred.

**Part b. [2 points] Identify the correct answers from these MCQs**

What is the Big O run-time for searching a value in a hash-table with linear probing <b>a.</b> $O(1)$ <b>b.</b> $O(\log n)$ <b>c.</b> $O(n)$ <b>d.</b> $O(n \log n)$	c
One of these is not a correct answer for reducing collisions in a hash table <b>a.</b> Use Linear probing method <b>b.</b> Use random values for hashing <b>c.</b> Use Separate Chaining method <b>d.</b> Use uniform hashing method	b
This data structure is useful for separate chaining <b>a.</b> Stacks <b>b.</b> Queue <b>c.</b> Linked Lists <b>d.</b> Heaps	c
A poor hash function would result in a hash table with: <b>a.</b> Clustering <b>b.</b> Uniform hashing <b>c.</b> No collisions <b>d.</b> No displacements	a

**Question 4. [4 + 4 = 8 points] [CLO 3]**

**Part a. [4 points]** Show the trace how the **Mergesort** algorithm sorts the given array, Identify the mid points for each pass/iteration.

K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S

**Part b. [4 points]** Give the correct answer for the following MCQs

<p>Given the following sequence: {2, 3, 5, 6, 9, 11, 15}. Which sorting algorithm will run in least time (n comparisons)?</p> <p>a. Insertion Sort  b. Selection Sort  c. Heap Sort  d. Merge Sort</p>	<b>a</b>
<p>Number of swaps/exchanges occurred sorting the data {0, 1, 2, 0, 1, 2, 0, 1, 2} using Selection sort.</p> <p>a. O(1)  b. O(log n)  c. O(n)  d. O(n log n)</p>	<b>c</b>
<p>Number of swaps/exchanges occurred sorting the data {0, 0, 0, 1, 1, 1, 2, 2, 2} using Insertion sort.</p> <p>a. O(1)  b. O(log n)  c. O(n)  d. O(n log n)</p>	<b>d</b>
<p>Worst-case run-time for quick-sort</p> <p>a. O(n<sup>2</sup>)  b. O(log n)  c. O(n<sup>3</sup>)  d. O(n log n)</p>	<b>A</b>
<p>Worst-case run-time for bottom-up merge-sort</p> <p>a. O(n<sup>2</sup>)  b. O(log n)  c. O(n<sup>3</sup>)  d. O(n log n)</p>	<b>D</b>
<p>Which if the following sorting algorithms requires extra space (i.e. not in-place)</p> <p>a. Insertion Sort  b. Selection Sort  c. Quick Sort  d. Merge Sort</p>	<b>d</b>
<p>Running mergesort on an array that is already sorted takes this time:</p> <p>a. O(1)  b. O(log n)  c. O(n)  d. O(n log n)</p>	<b>D</b>
<p>Which of the following algorithm design techniques is used in the quick sort algorithm</p> <p>a. Backtracking  b. Divide and Conquer  c. Uniform hashing  d. None of the above</p>	<b>B</b>

<This sheet is left blank intentionally. DO NOT detach>

<End of Exam>